Formal Verification of SciTokens



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NCSA | NATIONAL CENTER FOR SUPERCOMPUTING APPLICATIONS

Authentication bugs in SciToken will have catastrophic consequences to NSF projects



Our initial approach is to understand specs, code, and manual auditing.

SciToken Claims and Scopes

Language

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Standard Claims

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ciToken Claim Semantics

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SciTokens specs are well-defined



SciTokens code is sophisticated (~5K LOC)

These bugs allow unauthorized access of protected data. However, they are independent of SciToken's model correctness.

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Despite that SciToken has manual auditing, quality assurance (QA), and testing, they remain inadequate in unearthing critical bugs.



Formal verification of SciToken implementations is urgently needed.

Overview of Dafny: a verification-ready programming language



Putting Dafny in perspective



Microsoft Research

https://www.microsoft.com/en-us/research/project/boogie-an-intermediate-verification-language/

Why Dafny for scitokens?

- Static verification of programs at compile time, avoiding *data leak* and *system compromise* at runtime.
- Dafny compiler produces both the proof and cross-platform, verified executable code
 - Dafny code is compiled to .NET Common Intermediate Language (CIL)
 - CIL is translated into *six* languages including *compilea* (Java, Go, C++) and *interpretea* (Javascript, PHP, **Python**)



Example: verifying scitokens **audience**, a critical function, with Dafny.



Specs in natural language

Truth table

Python implementation

Errors & Counter examples

Example (2): verifying scitokens **scope** with Dafny.



scitokens.py



```
method Validate_Scope(
    value : string,
    test_access : bool,
    norm_requested_path : string,
    token scope : seq<seq<string>>)
    returns (
        t : bool,
        result : seq<seq<string>>)
requires |value| > 0;
    var scope := split1(value, ' ');
    var iter := 0;
    var authz := "":
    var norm_path := """;
    if (test_access == true) {
        while (iter < !scope!)
       invariant 0 <= iter <= |scope|;
            authz, norm_path, j := Check_Scope(scope[iter]);
```

Validate-scope.dfy (excerpt)

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Summary & Future Work

SciTokens Codebase & Specs

SciToken Claims and Scopes

Language

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Formal Verification

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Found two bugs





Path traversal attack

Verified three critical functions





Validate audience

Validate scope



Check scope

Correct-by-construction program synthesis

Learned formal verification tools

Hoare logic calculus $\{P\}C\{Q\}$

SMT/SAT solvers (z3, dafny)

Generated correct implementation in six languages

Compiled (Java, C++, Go)

Interpreted (PHP, Javascript, **Python**)

Putting human-in-the-loop with formal verification. 10